

Instrumentation in the era of Stack Map Frames

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About me

- Studies in program analysis and verification
- 5 years at Coverity (acquired by Synopsys)
- Technical lead for Java analysis

About Coverity

- “Development testing”
 - Defects found with static analysis
 - “Quality” and “security”
 - Defects found with dynamic analysis
 - Testing gaps found with test analysis
 - Build and architecture analysis
- www.coverity.com
- <http://scan.coverity.com>
- <http://code-spotter.com>

About Dynamic Analyzer

- Finds:
 - Resource leaks
 - Deadlocks
 - Race conditions
 - ... that might be too tricky for static analysis.
- Works via a Java agent doing bytecode instrumentation
- New in 7.6.0:
 - Agent injection
 - `cov-capture test command`
 - Java 8 support

About Dynamic Analyzer

- Finds:
 - Resource leaks
 - Deadlocks
 - Race conditions
 - ...
- Works
- New i
 - Age
 - c
 - Jav

harold seigel harold.seigel at oracle.com
Fri Mar 22 13:28:21 PDT 2013

Hi,

We have filed a CCC to deprecate -XX:-UseSplitVerifier in JDK 8 and it has been approved.

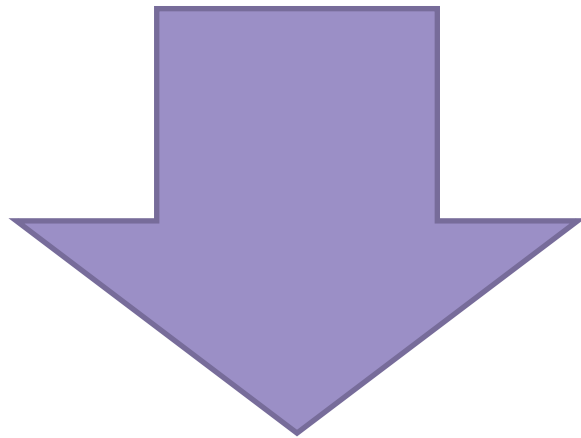
Previous users of -XX:-UseSplitVerify can initially use -noverify until they have correct version 51 bytecodes. Also, due to requests from folks who do generate version 51 bytecodes, we have added improvements to the new verifier error diagnostics - and specifically targeted them to give additional detailed information about stackmap tables. This should help in the transition to generating correct version 51 class files.

Thanks, Harold

The Story of Stack Maps

Thanks to
Prashant Deva
blog post

*Bytecode verification
is “too slow”*



*Bytecode requires types
on operand stack*

Java 5

Java 6 – experimental
feature

Java 7 – default feature,
with `-XX:-UseSplitVerifier`
work-around

Java 8 – No work-around.
Yet, JVM has old verifier!

Example

```
Comparator<String> getComparator(boolean b) {  
    return b ? new C1() : new C2();  
}
```

```
public java.util.Comparator<java.lang.String> get(boolean);
```

Code:

```
0: iload_1  
1: ifeq          14  
4: new          #18      // class C1  
7: dup  
8: invokespecial #20      // Method C1."<init>":()V  
11: goto         21  
14: new          #21      // class C2  
17: dup  
18: invokespecial #23      // Method C2."<init>":()V  
21: areturn
```

The diagram illustrates the control flow of the bytecode. A blue arrow points from instruction 1 (ifeq) to instruction 4 (new), indicating a branch to the C1 class initialization path. Another blue arrow points from instruction 11 (goto) to instruction 18 (invokespecial), indicating a branch to the C2 class initialization path.

The big problem

- How ASM deals with it:
 - Option 1: Define stack maps yourself
 - Option 2: Let ASM compute them
- BUT... instrumentation agent
 - Referenced classes might not be loaded yet

Solutions

- Don't inject new control flow
 - Just insert straight line code
- Consider adding indirection / helper methods
- Implement a super smart instrumentation API

The background features a stylized globe with various colored overlays and light trails. The globe is rendered in shades of blue and purple, with white outlines for continents. Overlaid on the globe are several glowing lines and patterns in red, green, and yellow, suggesting data flow or network connections. The overall aesthetic is futuristic and technological.

Thanks!

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www.coverity.com